



**Computer Science and Information Engineering
National Chi Nan University**

The Principle and Application of Secret Sharing

Dr. Justie Su-Tzu Juan

Lecture 5. Perfect SSS for Graph-Based Structure

§ 5.1 Perfect Secret Sharing Schemes

**Slides for a Course Based on
Y.-F. Weng “A Study of Perfect Secret Sharing Scheme”, Master Thesis of
Department of SCIE, National Chi Nan University, 2006.**



§ 5.1 Perfect Secret Sharing Schemes

- **Def:**

- Let \mathcal{K} be the master key space and \mathcal{S}_i be the share space for participant i . The *information rate* ρ of the secret sharing scheme is defined as $\rho = \min_i \log_2 |\mathcal{K}| / \log_2 |\mathcal{S}_i|$.
- A secret sharing scheme is *ideal* if $\rho = 1$.
- The *minimal access structure* $\Gamma_0 = \{A \in \Gamma : A' \not\subseteq A \text{ for all } A' \in \Gamma - \{A\}\}$.
- The *maximal prohibited structure* $\Delta_1 = \{B \in \Delta : B \not\subseteq B' \text{ for all } B' \in \Gamma - \{B\}\}$.

$P = \{P_1, P_2, P_3\}$
 $\Gamma = \{\{P_1, P_2\}, \{P_1, P_3\}, \{P_1, P_2, P_3\}\}$
 $\Delta = \{\{P_1\}, \{P_2\}, \{P_3\}, \{P_1, P_2\}\}$

Γ_0 (red box around $\{P_1, P_2\}$)
 Δ_0 (blue box around $\{P_3\}$)



§ 5.1 Perfect Secret Sharing Schemes

- **Def:**

- **Graph** $G = (V, E)$

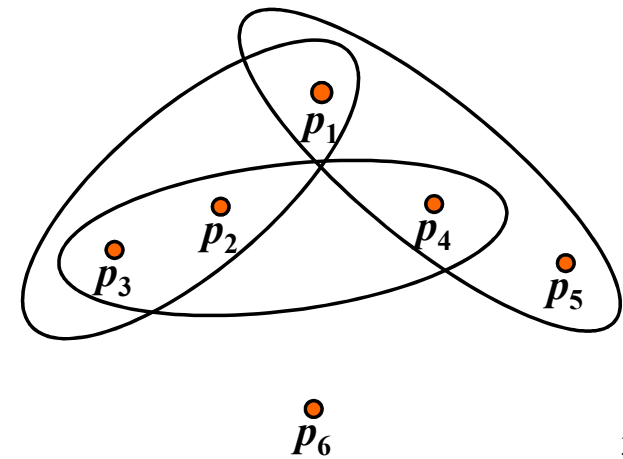
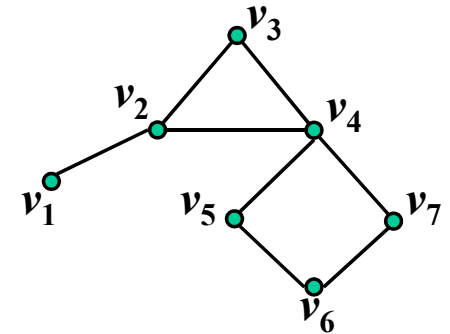
- $E = \{e_1, e_2, \dots, e_\varepsilon\}$ and $e_i = \{u, v\}$ where $u, v \in V, 1 \leq i \leq \varepsilon$

- **Hypergraph** $H = (V, E)$

- $E = \{E_1, E_2, \dots, E_{|E|}\}$ and $|E_i| \geq 2, 1 \leq i \leq |E|$

- **r -uniform Hypergraph:**

- $\forall |E_i| = r, 1 \leq i \leq |E|$





§ 5.1 Perfect Secret Sharing Schemes

- **Related works:**
- Perfect SSS for **graph-based prohibited structure (Type II)**
 - SS scheme(1997)
 - Sun, Shieh
 - Sun's scheme (1999)
- Perfect SSS for **general access structure**
 - Tochikubo's scheme (2004) (Type I)
 - TUM scheme (2005) (Type II)
 - Tochikubo, Uyematsu, Matsumoto



§ 5.1 Perfect Secret Sharing Schemes

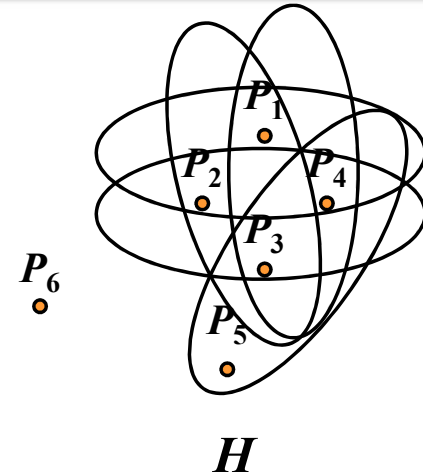
- r -uniform hypergraph-based prohibited structure**

- r -uniform hypergraph $H = (V, E)$

- $V(H) = P$ and $|P| = n$

- $\Delta = \{A: A \subseteq P \text{ and } |A| \leq r - 1\} \cup E(H)$

- $\Gamma = \{A: A \subseteq P \text{ and } |A| \geq r + 1\} \cup \{A: A \notin E(H) \text{ and } |A| = r\}$



- Weng and Juan?

$$P = \{p_1, p_2, p_3, p_4, p_5, p_6\}$$

- Weng and Juan

$$\Delta = \{A: A \subseteq P \text{ and } |A| \leq 2\} \cup E(H)$$

- Weng, Juan and

$$= \{\emptyset, \{p_1\}, \{p_2\}, \{p_3\}, \{p_4\}, \{p_5\}, \{p_6\}, \{p_1, p_2\}, \{p_1, p_3\}, \{p_1, p_4\}, \{p_1, p_5\}, \{p_1, p_6\}, \{p_2, p_3\}, \{p_2, p_4\}, \{p_2, p_5\}, \{p_2, p_6\}, \{p_3, p_4\}, \{p_3, p_5\}, \{p_3, p_6\}, \{p_4, p_5\}, \{p_4, p_6\}, \{p_5, p_6\}, \{p_1, p_2, p_3\}, \{p_1, p_2, p_4\}, \{p_1, p_3, p_4\}, \{p_2, p_3, p_4\}, \{p_3, p_4, p_5\}\}$$

$$\Gamma = \{A: A \subseteq P \text{ and } |A| \geq 4\} \cup \{\{p_1, p_2, p_5\}, \{p_1, p_2, p_6\}, \{p_1, p_3, p_5\}, \{p_1, p_3, p_6\}, \{p_1, p_4, p_5\}, \{p_1, p_4, p_6\}, \{p_1, p_5, p_6\}, \{p_2, p_3, p_5\}, \{p_2, p_3, p_6\}, \{p_2, p_4, p_5\}, \{p_2, p_4, p_6\}, \{p_2, p_5, p_6\}, \{p_3, p_4, p_6\}, \{p_3, p_5, p_6\}, \{p_4, p_5, p_6\}\}$$



§ 5.1 Perfect Secret Sharing Schemes

- r -uniform hypergraph-based prohibited structure

	Sun-Shieh (1997)	Sun (1999)	Tochikubo (2004)	TUM (2005)
extend	r -HP1 3-1	r -HP2 4-1		I-TUM 6-2
VD	r -VDHP1 3-2	r -VDHP2 4-2	VDT 5-2	VDTUM
M	r -MHP1 3-3	r -MHP2 4-3	MT 5-3	MITUM 6-3
VDM	r -VDMHP1 3-4	r -VDMHP2 4-4	VDMT 5-4	VDMITUM 6-4



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Lecture 5. Perfect SSS for Graph-Based Structure

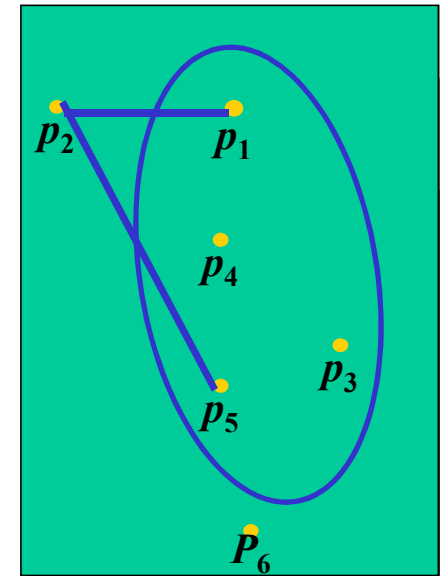
§ 5.2 Hypergraph-based SSS for General Access Structures

Slides for a Course Based on
Y.-C. Wang “Using Hypergraph to Design Perfect Secret Sharing Schemes
for General Access Structures”, Master Thesis of Department of SCIE,
National Chi Nan University, 2007.



§ 5.2 Hypergraph-based SSS for General Access Structures

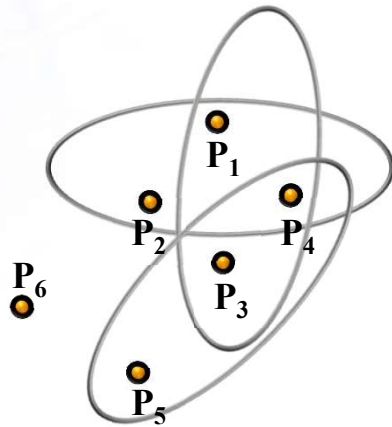
- **Hypergraph**
 - r -uniform hypergraph
 - $r = 2$: graph
 - $r > 2$
 - (r_1, r_2) -uniform hypergraph
 - (r_1, r_2, r_3) -uniform hypergraph
 - General hypergraph
- **Hypergraph-based Access structure**
 - $\Gamma = \{A \subseteq P : S \subseteq A \text{ for any } S \in \Delta_0\}$
 - $\Delta = 2^P \setminus \Gamma = \{A \subseteq P : S \not\subseteq A \text{ for all } S \in \Delta_0\}$.



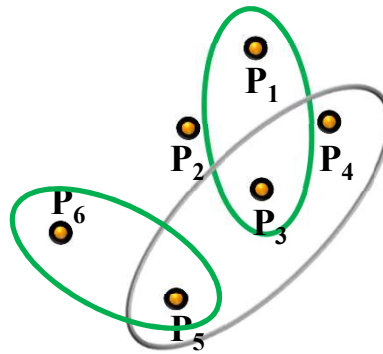


§ 5.2 Hypergraph-based SSS for General Access Structures

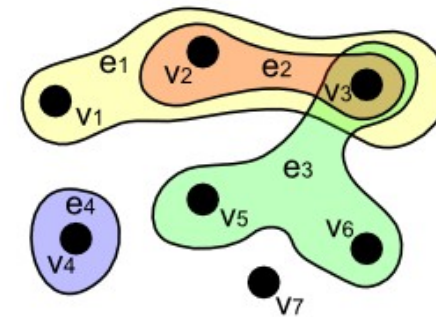
- Hypergraph $H = (V, E)$
 - r -Uniform Hypergraph
 - (r_1, r_2) -Uniform Hypergraph
 - General Hypergraph



3-Uniform Hypergraph



(2, 3)-Uniform Hypergraph



General Hypergraph

Source: Wikipedia



§ 5.2 Hypergraph-based SSS for General Access Structures

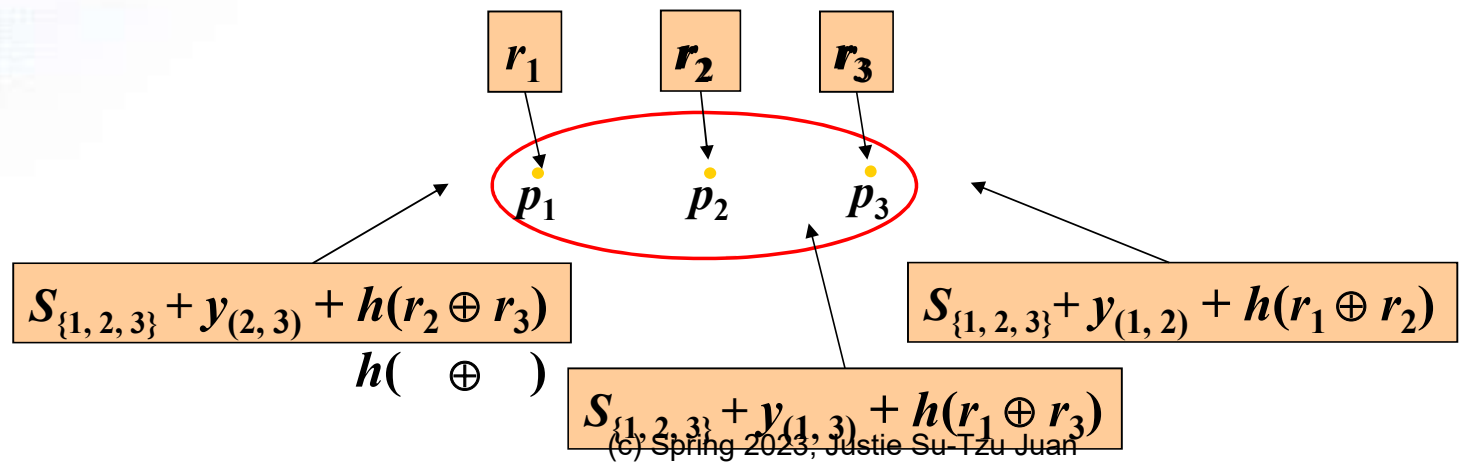
- r -HA Scheme (2007)
Structures
- (r_1, r_2) -HA Scheme (2007)
Access Structures
- (r_1, r_2, r_3) -HA Scheme
Based Access Structures
- G-HA Scheme
 - G-VDHA Scheme
 - G-MHA Scheme
 - G-VDMHA Scheme



§ 5.2 Hypergraph-based SSS for General Access Structures

- ***r*-HA Scheme - Idea**
- 3-Uniform Hypergraph
 - $K = (K_1, K_0) \in Z_q \times Z_q$
 - $f(x) = k_2x^2 + K_1x + K_0 \pmod{q}$
 - $y_{(i,j)} = f(i \cdot n + j)$
 - h : a one-way hash function

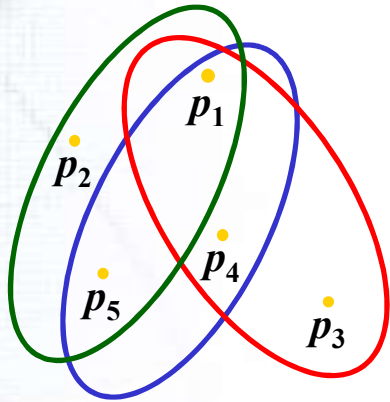
$$\begin{aligned}
 S_{\{1,2,3\}} = & y_{(1,2)} + h(r_1 \oplus r_2) \\
 & + y_{(2,3)} + h(r_2 \oplus r_3) \\
 & + y_{(1,3)} + h(r_1 \oplus r_3)
 \end{aligned}$$





§ 5.2 Hypergraph-based SSS for General Access Structures

• Ex:



S_1	S_2
r_1	r_2
$S_{\{1,2,5\}} + y_{(2,5)} + h(r_2 \oplus r_5)$	$S_{\{1,2,5\}} + y_{(1,5)} + h(r_1 \oplus r_5)$
$S_{\{1,3,5\}} + y_{(3,4)} + h(r_3 \oplus r_4)$	—
$S_{\{1,4,5\}} + y_{(4,5)} + h(r_4 \oplus r_5)$	—

S_3	S_4	S_5
r_3	r_4	r_5
—	—	$S_{\{1,2,5\}} + y_{(1,2)} + h(r_1 \oplus r_2)$
$S_{\{1,3,5\}} + y_{(1,4)} + h(r_1 \oplus r_4)$	$S_{\{1,3,5\}} + y_{(1,3)} + h(r_1 \oplus r_3)$	—
	$S_{\{1,4,5\}} + y_{(1,5)} + h(r_1 \oplus r_5)$	$S_{\{1,4,5\}} + y_{(1,4)} + h(r_1 \oplus r_4)$

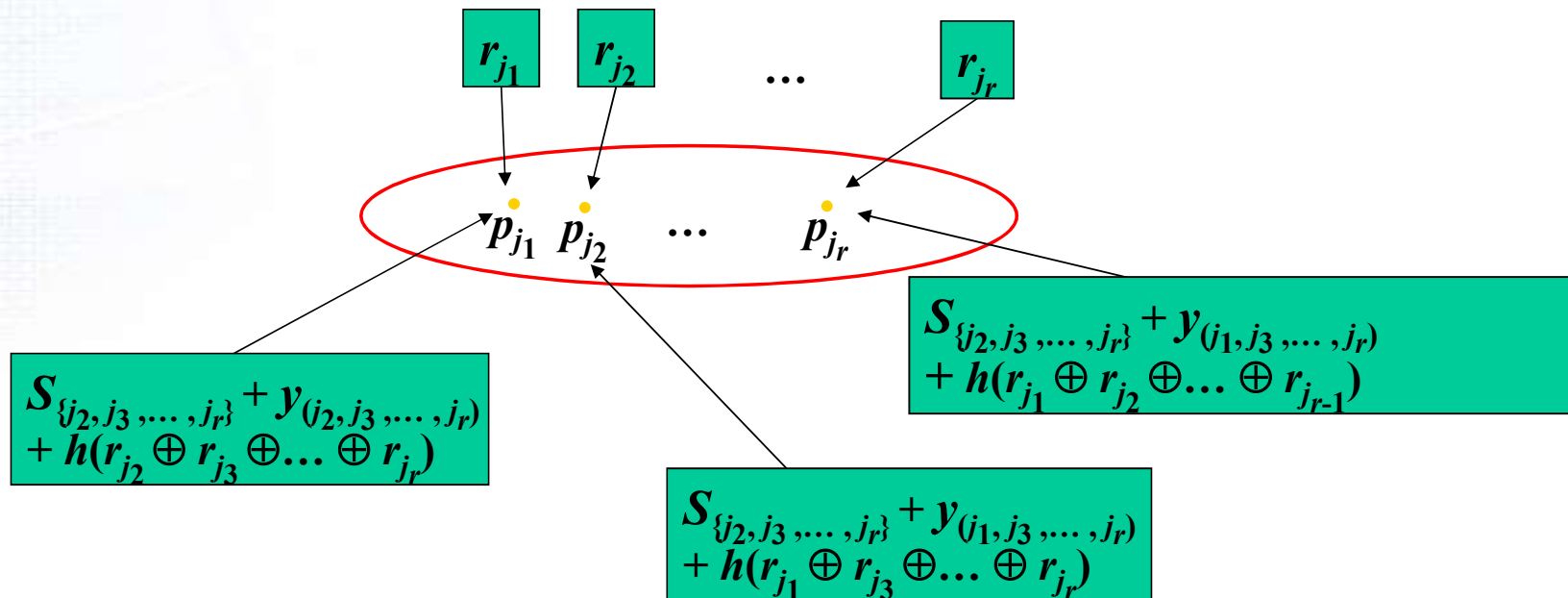
$$S_{\{1,2,5\}} = y_{(1,2)} + h(r_1 \oplus r_2) + y_{(2,5)} + h(r_2 \oplus r_5) + y_{(1,5)} + h(r_1 \oplus r_5)$$

...



§ 5.2 Hypergraph-based SSS for General Access Structures

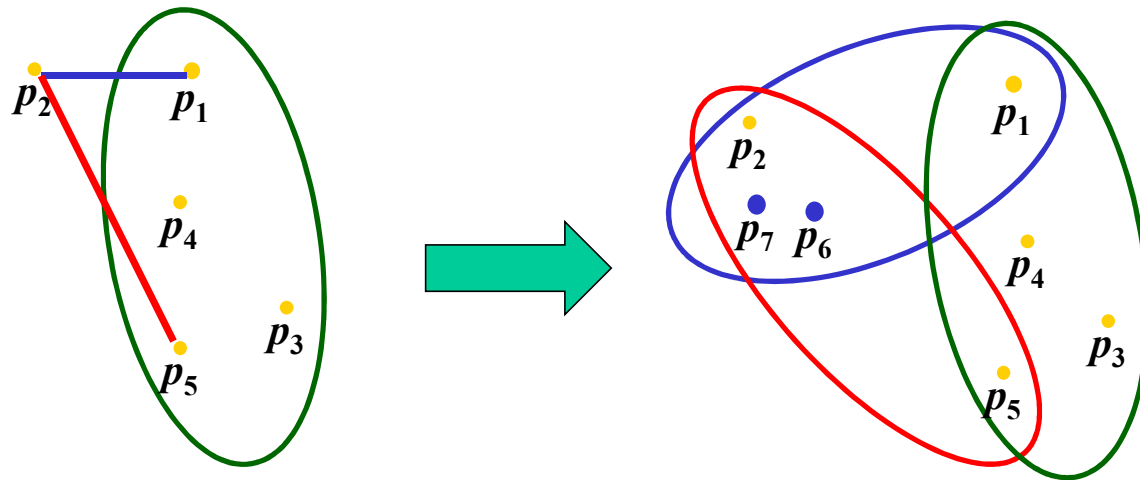
- $f(x) = k_{r-1}x^{r-1} + \dots + k_2x^2 + K_1x + K_0 \pmod{q}$
- $y_{(i_1, i_2, \dots, i_{r-1})} = f\left(\sum_{i=1}^{r-1} i_k n^{r-k-1}\right)$





§ 5.2 Hypergraph-based SSS for General Access Structures

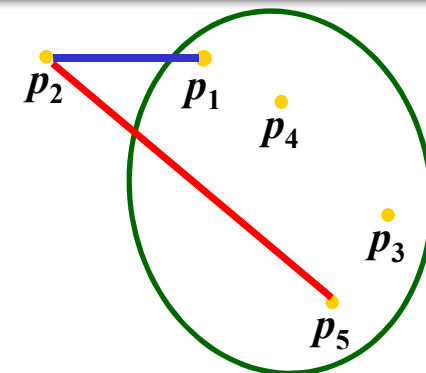
- (r_1, r_2) -HA Scheme – Idea
 - (2, 4)-HA Scheme



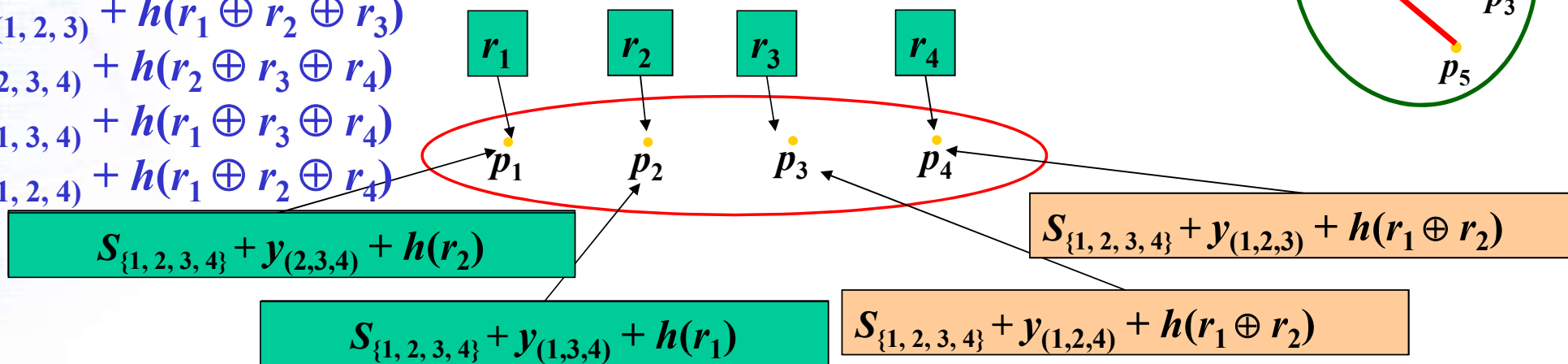


§ 5.2 Hypergraph-based SSS for General Access Structures

- (r_1, r_2) -HA Scheme
 - (2, 4)-HA Scheme



$$\begin{aligned}
 S_{\{1,2,3,4\}} &= y_{(1,2,3)} + h(r_1 \oplus r_2 \oplus r_3) \\
 &+ y_{(2,3,4)} + h(r_2 \oplus r_3 \oplus r_4) \\
 &+ y_{(1,3,4)} + h(r_1 \oplus r_3 \oplus r_4) \\
 &+ y_{(1,2,4)} + h(r_1 \oplus r_2 \oplus r_4)
 \end{aligned}$$



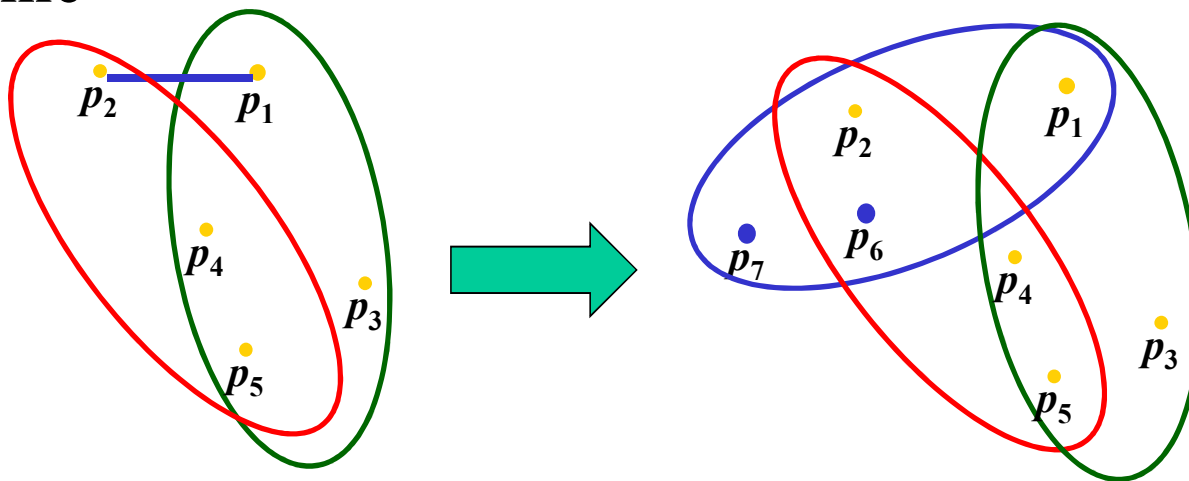
S_1	S_2	S_3	S_4	S_5	S_6	S_7
r_1	r_2	r_3	r_4	r_5	0	0
$S_{\{1,2,3,4\}} + y_{(3,4,5)} + h(r_3 \oplus r_4 \oplus r_5)$	-	$S_{\{1,2,3,4\}} + y_{(1,4,5)} + h()$	$S_{\{1,2,3,4\}} + y_{(1,3,5)} + h()$	$S_{\{1,2,3,4\}} + y_{(1,3,4)} + h()$		
$S_{\{1,2,6,7\}} + y_{(2,6,7)} + h(r_2)$	$S_{\{1,2,6,7\}} + y_{(1,6,7)} + h(r_1)$	-	-	-	$S_{\{1,2,6,7\}} + y_{(1,2,7)} + h(r_1 \oplus r_2)$	$S_{\{1,2,6,7\}} + y_{(1,2,6)} + h(r_1 \oplus r_2)$
-	$S_{\{2,5,6,7\}} + y_{(5,6,7)} + h(r_5)$	-	(c) Spring 2023, Justice Su-Tzu Juan	$S_{\{2,5,6,7\}} + y_{(2,6,7)} + h(r_2)$	$S_{\{2,5,6,7\}} + y_{(2,5,7)} + h(r_2 \oplus r_5)$	$S_{\{2,5,6,7\}} + y_{(2,5,6)} + h(r_2 \oplus r_5)$



§ 5.2 Hypergraph-based SSS for General Access Structures

- (r_1, r_2, r_3) -HA Scheme

- (2, 3, 4)-HA Scheme



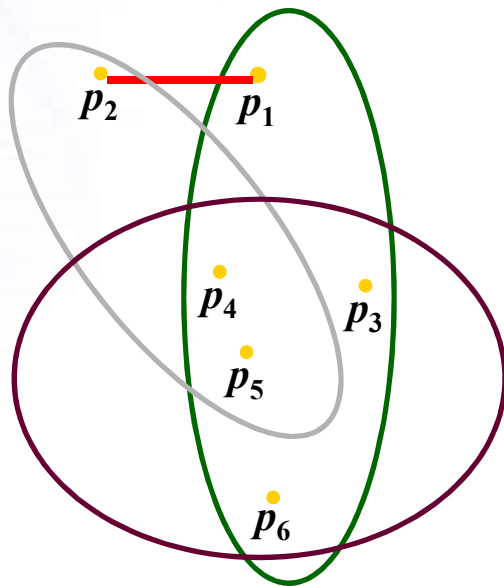
$$\begin{aligned}
 S_{\{1,3,4,5\}} &= y_{(1,3,4)} + h(r_1 \oplus r_3 \oplus r_4) \\
 &+ y_{(1,3,5)} + h(r_1 \oplus r_3 \oplus r_5) \\
 &+ y_{(1,4,5)} + h(r_1 \oplus r_4 \oplus r_5) \\
 &+ y_{(3,4,5)} + h(r_3 \oplus r_4 \oplus r_5)
 \end{aligned}$$

S_1	S_2	S_3	S_4	S_5	S_6	S_7
r_1	r_2	r_3	r_4	r_5	0	0
$S_{\{1,3,4,5\}} + y_{(3,4,5)} + h(r_3 \oplus r_4 \oplus r_5)$	—	$S_{\{1,3,4,5\}} + y_{(1,4,5)} + h$	$S_{\{1,3,4,5\}} + y_{(1,3,5)} + h$	$S_{\{1,3,4,5\}} + y_{(1,3,4)} + h$	—	—
—	$S_{\{2,4,5,6\}} + y_{(4,5,6)} + h(r_4 \oplus r_5)$	—	$S_{\{2,4,5,6\}} + y_{(2,5,6)} + h$	$S_{\{2,4,5,6\}} + y_{(2,4,6)} + h$	$S_{\{2,4,5,6\}} + y_{(2,4,5)} + h$	—
$S_{\{1,2,6,7\}} + y_{(2,6,7)} + h(r_2)$	$S_{\{1,2,6,7\}} + y_{(1,6,7)} + h(r_1)$	—	—	—	$S_{\{1,2,6,7\}} + y_{(1,2,7)} + h$	$S_{\{1,2,6,7\}} + y_{(1,2,6)} + h$

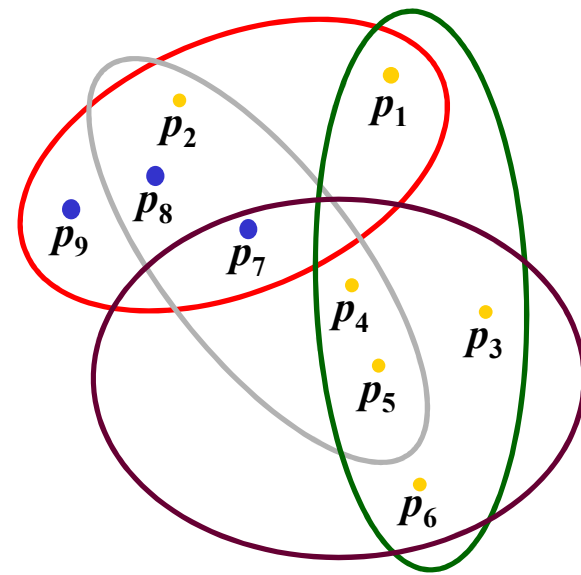


§ 5.2 Hypergraph-based SSS for General Access Structures

- **G-HA Scheme – Idea**
 - $(r_1, r_2, \dots, r_\omega)$ -HA Scheme
 - $(2, 3, 4, 5)$ -HA Scheme



$$|V| = n$$

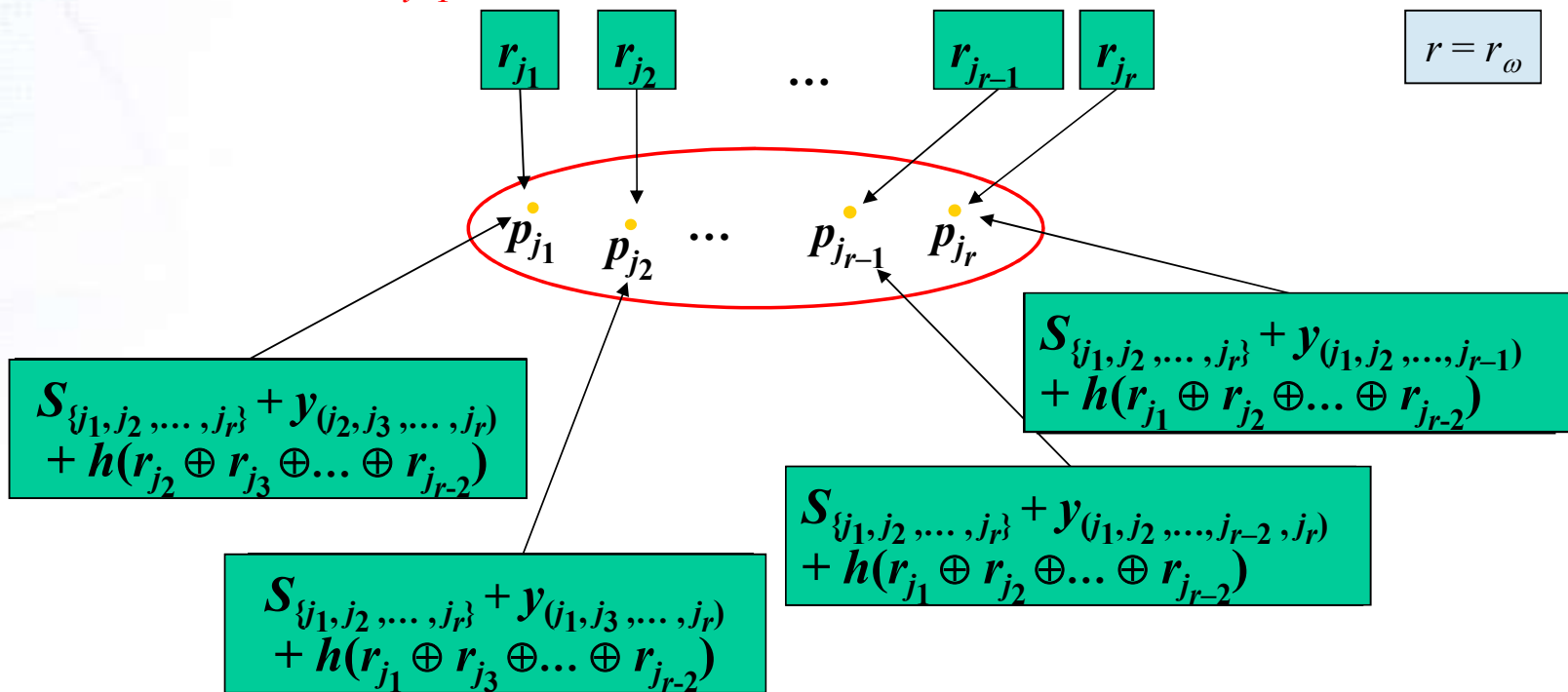


$$|V| = n'$$



§ 5.2 Hypergraph-based SSS for General Access Structures

- $f(x) = k_{r-1}x^{r-1} + \dots + k_2x^2 + K_1x + K_0 \pmod{q}$
- $y_{(i_1, i_2, \dots, i_{r-1})} = f\left(\sum_{i=1}^{r-1} i_k(n^i)^{r-k-1}\right)$





§ 5.2 Hypergraph-based SSS for General Access Structures

- **G-VDHA Scheme**

Publish:

$$g, g^{K_0}, g^{K_1}, g^{k_2} \text{ mod } p, g^{S_{\{1,2,3\}}}$$

$$g^{r_1}, g^{r_2}, g^{h(r_1)}, g^{h(r_2)} \text{ mod } p$$

g : generator

To Check p_2 :

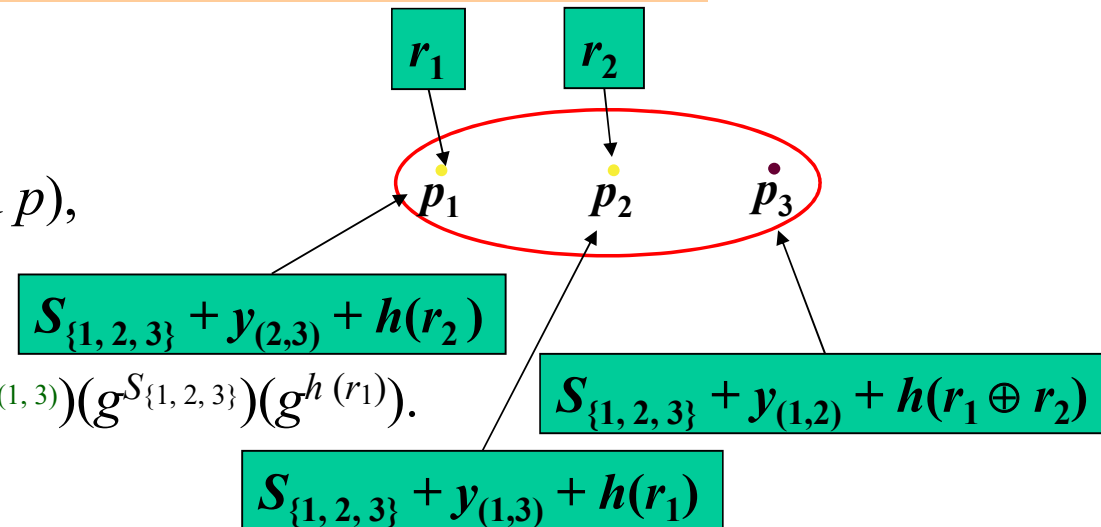
– Calculate

$$g^{y(1,3)} = (g^{k_2})^{x^2} (g^{K_1})^{x^1} (g^{K_0}) \text{ (mod } p),$$

where $x = 1 \cdot 3 + 3$.

– Check $g^{y(1,3)+S_{\{1,2,3\}}+h(r_1)} = (g^{y(1,3)})(g^{S_{\{1,2,3\}}})(g^{h(r_1)})$.

– Check $g^{r_2} = (g)^{r_2}$.





§ 5.2 Hypergraph-based SSS for General Access Structures

- **G-MHA Scheme**

Publish:

F : one way hash function

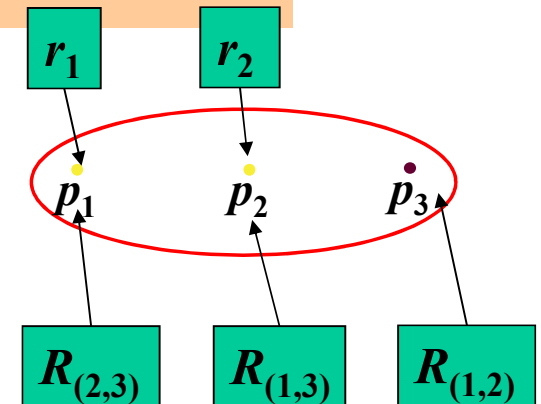
random numbers R over Z_q

$$y_{(1,2)} + S_{\{1,2,3\}} + h(F(r_1, R) \oplus F(r_2, R)) + F(R_{(1,3)}, R)$$

$$y_{(1,3)} + S_{\{1,2,3\}} + h(F(r_1, R)) + F(R_{(1,3)}, R)$$

$$y_{(2,3)} + S_{\{1,2,3\}} + h(F(r_2, R)) + F(R_{(2,3)}, R)$$

- If $\{p_1, p_3, p_3\}$ want to reconstruct the K :
 - $p_1: F(R_{(2,3)}, R), F(r_1, R);$
 - $p_2: F(R_{(1,3)}, R), F(r_2, R);$





§ 5.2 Hypergraph-based SSS for General Access Structures

- **G-VDMHA Scheme**

Publish:

random numbers R over Z_q

$$y_{(1,2)} + h(F(r_1, R) \oplus F(r_2, R)) + F(R_{(1,2)}, R)$$

$$y_{(1,3)} + h(F(r_1, R)) + F(R_{(1,3)}, R)$$

$$y_{(2,3)} + h(F(r_2, R)) + F(R_{(2,3)}, R)$$

$$g, g^{K_0}, g^{K_1}, g^{K_2}, g^{F(r_1, R)}, g^{F(r_2, R)}, g^{F(r_3, R)},$$

$$g^{h(F(r_1, R) \oplus F(r_2, R))}, g^{h(F(r_1, R))}, g^{h(F(r_2, R))}.$$

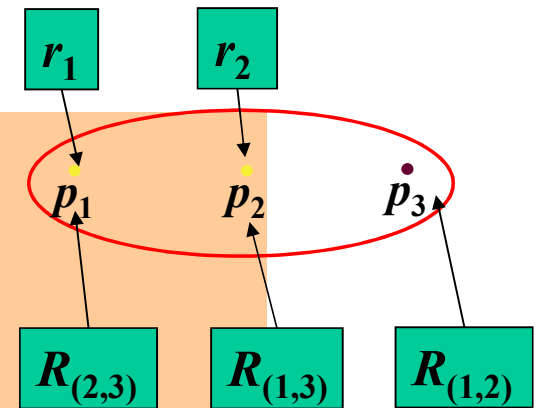
Calculate

$$g^{y(1,3)} = (g^{K_2})^{x^2} (g^{K_1})^{x^1} (g^{K_0}) \pmod{p}, \text{ where } x = 1 \cdot n' + 3.$$

Check

$$g^{y(1,3) + h(F(r_1, R)) + F(R_{(1,3)}, R)} = (g^{y(1,3)}) (g^{h(F(r_1, R))}) (g^{F(R_{(1,3)}, R)})$$

$$\text{Check } g^{F(r_1, R)} = (g)^{F(r_1, R)}.$$



$$p_1: F(R_{(2,3)}, R), F(r_1, R);$$

$$p_2: F(R_{(1,3)}, R), F(r_2, R).$$



§ 5.2 Hypergraph-based SSS for General Access Structures

$d =$ the minimum degree of G ;
 $m_i = \{A: A \in \Gamma_0, |A| = r_i\}$;
 $m = \sum_{i=1}^{\omega} m_i$;
 $c_i = \max_{A \in \Gamma_0} |A| - r_i$.

- Performance**

	G-HA scheme (2007)	Tsai et al.'s scheme (1999)	Wang's scheme (2004)
Information rate	$2 / (d + 1)$	1	1
multiple operation	$m \times r_{\omega}^2 = \sum_{p_i} d_i \times m_i$	$r_i \times m_i$	$\sum_{p_i} (d_i + 2)$
addition operation	$m \times r_{\omega}^2 = \sum_{p_i} d_i \times m_i$	0	0
power operation	0	$2n + 1$	n
hash function	$\sum_{p_i} d_i$	0	0
exclusive-or	$\sum_{p_i} d_i \times (r_{\omega} - 1)$	m_i	0
inverse	0	0	$n + m$
Gaussian elimination	0	0	1
#Pseudo-man	$(r_{\omega} - r_1) / 0$	0	0



§ 5.2 Hypergraph-based SSS for General Access Structures

$d =$ the minimum degree of G ;
 $m_i = \{A: A \in \Gamma_0, |A| = r_i\}$;
 $m = \sum_{i=1}^{\omega} m_i$;
 $c_i = \max_{A \in \Gamma_0} |A| - r_i$.

- Performance**

	TUM Scheme* (2005)	G-HA scheme
information rate	$1 / d$	$2 / (d + 1)$
the number of public share	0	$\sum_{i=1}^{\omega} (m_i \times c_i)$
space complexity	$\sum_{i=1}^{\omega} m_i (r_i - 1) = O(mr_{\omega})$	$r_{\omega} - 1 = O(r_{\omega})$
multiple operation	$\sum_{i=1}^{\omega} (m_i \times r_i^2) = O(mr_{\omega}^2)$	$m \times r_{\omega}^2 = O(mr_{\omega}^2)$
addition operation	$\sum_{i=1}^{\omega} (m_i \times r_i^2) = O(mr_{\omega}^2)$	$m \times r_{\omega}^2 = O(mr_{\omega}^2)$
exclusive-or operation	0	$m r_{\omega} (r_{\omega} - 1) = O(mr_{\omega}^2)$
hash function	0	$(r_{\omega} - 1) \times m = O(mr_{\omega})$

* The hypergraph is proper.



§ 5.2 Hypergraph-based SSS for General Access Structures

- **Modification:**

- $K = (K_1, K_0) \in Z_q \times Z_q$
 $f(x) = k_{r-1}x^{r-1} + \dots + k_2x^2 + K_1x + K_0 \pmod{q}$

- $K = (K_r, \dots, K_2, K_1, K_0) \in [Z_q]^r$
 $f(x) = K_r x^{r-1} + \dots + K_2 x^2 + K_1 x + K_0 \pmod{q}$

Non-perfect !