



# Graph Theory (圖型理論)

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# Introduction

- 圖論為一在二十世紀快速成長的一門學問。自原本只是組合數學書中一個章節的短短數頁，百倍增長為如今此一可單獨研討的學門。它之所以會如此快速成長，主要的原因是圖論的應用性非常之廣。諸如：物理、化學、生物、心理學、社會學等等。而另一個原因則是，在電腦科學理論上，一些關於處理複雜度的問題，能被轉換成圖論問題來解決。



# Introduction

- 師生晤談時間：二EF
- 主要教科書：Junming Xu, Theory and Application of Graphs, Kluwer Academic Publishers (科大文化代理), 2003.
- 重要參考書籍：
  1. West, Introduction to Graph Theory 2ed, Prentice-Hill (全華代理), 2001.
  2. Bondy and Murty, Graph Theory with Application. 1976.
  3. Balakrishnan and Ranganatha, A Textbook of Graph Theory 2ed, Springer, 2012.
  4. 張鎮華, 演算法觀點的圖論, 台灣大學出版社, 2017.



# Introduction

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- 課程內容：預計將介紹下列各項
  1. Basic Concepts of Graphs
  2. Trees and Graphical Spaces
  3. Planar Graphs and Planar Graphs
  4. Flows and Connectivity
  5. Matchings and Independent Sets
  6. Coloring Theory



# Introduction

- 評量方式：最高分99  
作業40% + 平時成績20% + 期中考20% + 期末報告 30% + 加分作業  
大學部及格60, 研究生及格70
- 進度：  
第七週、第十週參加會議，補課。  
11/29 期中考：Chap 1 ~ Chap 2.  
1/3, 1/10：期末報告
- 助教：王怡君：計算理論研究室R307-1 (分機4862)  
Office Hour：無，但可以隨時傳訊息問問題。
- 網頁：<http://www.csie.ncnu.edu.tw/~jsjuan/courses.html>



# 圖型理論課程 核心能力 與 課程地圖



# 暨大資工系 教育目標

研發潛能  
理論能力

產業需求  
實作能力

人的本質  
自己以外

## 研究所教育目標

1. 配合國家經濟發展，培養符合資訊產業需求的工程技術人才
2. 配合國家科技及學術發展，培養具備前瞻資訊科技研發能力的人才
3. 配合全球永續發展潮流，培養具備國際視野、工程倫理、人文關懷及社會責任的科技人才



# 暨大資工系 核心能力

基礎數理  
理論

程式設計

論文撰寫

研究所核心能力

1. 具備資訊科學基礎數理知識並應用於發掘、分析與解釋數據的能力
2. 具備程式設計基礎知識並應用於設計及實作資訊軟體的能力
3. 具備使用英文閱讀資訊領域技術文件及學術論文的能力
4. 具備團隊合作及獨立執行資訊工程領域學術研究的能力
5. 具備撰寫學術論文的能力
6. 理解資訊工程專業倫理、敬業態度、環境保護及社會責任

英文能力

合作與獨立

生命品格





# 暨大資工系 課程地圖 (部分)

暨大資工系課程綜覽-100 學年度 (2011/09/13 修訂版)

基礎課程 (必修課程) 72 學分				專業領域 (選修) 32 學分
大一	大二	大三	大四	開課年級依各課程決定
普通物理 (上)	工程數學	微算機系統		多媒體領域
微積分 (上)	邏輯設計	系統程式		<b>演算法與計算理論領域</b>
離散數學	數位電子學	資料庫系統		數位網路通訊領域
計算機概論	資料結構與演算法 (一)	專題 (一)		訊號與資訊處理領域
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普通物理 (下)	資料結構與演算法 (二)	作業系統		硬體與電路領域
微積分 (下)	機率	編譯器		全校共同課程14學分
程式設計	計算機組織與結構	計算機網路		通識領域課程17學分
電子電路	線性代數	專題 (二)		專業選修 $\geq 32$ 學分
	數位電路實驗	微算機實驗		未來發展 (職涯)





# 暨大科技學院 核心能力

- ➔ ● 1. 專業知識與實務技能
- ➔ ● 2. 創新與獨立思考能力
- 3. 溝通表達與團隊合作精神
- ➔ ● 4. 專業倫理與社會責任認知
- 5. 掌握國際趨勢與全球視野



# 暨大學生 八大基本素養與核心能力

- ● (一) 道德思辨與實踐能力
- (二) 人際溝通與表達能力
- ● (三) 獨立思考與創新能力
- ● (四) 人文關懷與藝術涵養
- ● (五) 專業知能與數位能力
- (六) 團隊合作與樂業倫理
- (七) 全球視野與尊重多元文化
- (八) 社區參與與公民責任



# **Chapter 1**

# **Basic Concepts of Graphs**

## **§ 1.1 Graph and Graphical Presentation**



# 1.1 Graph and Graphical Presentation

- Def:
  - A **graph**  $G$  is an ordered triple  $(V, E, \psi)$ , where  
 $V, E$ : disjoint sets;  $\psi: E \rightarrow V \times V$ : a mapping.
  - $V$ : **vertex-set**;  $x \in V$ : **vertex**  
 $E$ : **edge-set**;  $e \in E$ : **edge**  
 $\psi$ : **incidence function**; if  $\psi(e) = (x, y)$ :  $x, y$  are **end-vertices** of  $e$ .
- Example 1.1.1:

$D = (V(D), E(D), \psi_D)$  is a digraph, where

$$V(D) = \{x_1, x_2, x_3, x_4, x_5\},$$
$$E(D) = \{a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9\}$$

and  $\psi_D$  is defined by  $\psi_D(a_1) = (x_1, x_2)$ ,  $\psi_D(a_2) = (x_3, x_2)$ ,  $\psi_D(a_3) = (x_3, x_3)$ ,  
 $\psi_D(a_4) = (x_4, x_3)$ ,  $\psi_D(a_5) = (x_4, x_2)$ ,  $\psi_D(a_6) = (x_4, x_2)$ ,  
 $\psi_D(a_7) = (x_5, x_2)$ ,  $\psi_D(a_8) = (x_2, x_5)$ ,  $\psi_D(a_9) = (x_3, x_5)$ .



# 1.1 Graph and Graphical Presentation

- Example 1.1.2:

$H = (V(H), E(H), \psi_H)$  is a digraph, where

$$V(H) = \{y_1, y_2, y_3, y_4, y_5\},$$

$$E(H) = \{b_1, b_2, b_3, b_4, b_5, b_6, b_7, b_8, b_9\}$$

and  $\psi_H$  is defined by  $\psi_H(b_1) = (y_1, y_2)$ ,  $\psi_H(b_2) = (y_3, y_2)$ ,  $\psi_H(b_3) = (y_3, y_3)$ ,

$$\psi_H(b_4) = (y_4, y_3), \psi_H(b_5) = (y_4, y_2), \psi_H(b_6) = (y_4, y_2),$$

$$\psi_H(b_7) = (y_5, y_2), \psi_H(b_8) = (y_2, y_5), \psi_H(b_9) = (y_3, y_5).$$

- Def:

- If  $V \times V$  is a set of ordered pair  $(x, y)$ 's, then

- ①  $G$  is called **directed graph (digraph)**.

- ②  $e \in E$ ; **directed edge (arc)**.

- ③ if  $\psi(e) = (x, y)$ :  $x$  is called the **tail** of  $e$ ;  $y$  is called the **head** of  $e$ ;

$e$  is called an **out-going edge** of  $x$ ; **in-coming edge** of  $y$ .



# 1.1 Graph and Graphical Presentation

- Def:
  - If  $V \times V$  is a set of unordered pair  $\{x, y\}$ 's, then
    - ①  $G$  is called an **undirected graph**.
    - ② Use  **$xy$**  or  **$yx$**  instead of  $\{x, y\}$ .
    - ③  $e \in E$ : **undirected edges**.
- Example 1.1.3:

$G = (V(G), E(G), \psi_G)$  is an undirected graph, where

$$V(G) = \{z_1, z_2, z_3, z_4, z_5, z_6\},$$
$$E(G) = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7, e_8, e_9\}$$

and  $\psi_G$  is defined by  $\psi_G(e_1) = z_1z_2$ ,  $\psi_G(e_2) = z_1z_4$ ,  $\psi_G(e_3) = z_1z_6$ ,  
 $\psi_G(e_4) = z_2z_3$ ,  $\psi_G(e_5) = z_3z_4$ ,  $\psi_G(e_6) = z_3z_6$ ,  
 $\psi_G(e_7) = z_2z_5$ ,  $\psi_G(e_8) = z_4z_5$ ,  $\psi_G(e_9) = z_5z_6$ .



# 1.1 Graph and Graphical Presentation

- Def:
  - **loop**: ex:  $\psi(e) = (x, x)$ .
  - **parallel edges (multi-edges)**: ex:  $\psi(e_1) = (x, y) = \psi(e_2)$
  - $E_G(x, y) = \{e \in E(G) : \psi(e) = (x, y)\}$
  - $\mu_G(x, y) = |E_G(x, y)|$
  - $\mu(G) = \max\{\mu_G(x, y) : \forall x, y \in V(G)\}$  : the **multiplicity** of  $G$ .

- ex: in Example 1.1.1:

$D = (V(D), E(D), \psi_D)$  is a digraph, where

$$V(D) = \{x_1, x_2, x_3, x_4, x_5\},$$

$$E(D) = \{a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9\}$$

and  $\psi_D$  is defined by  $\psi_D(a_1) = (x_1, x_2)$ ,  $\psi_D(a_2) = (x_3, x_2)$ ,  $\psi_D(a_3) = (x_3, x_3)$ ,  
 $\psi_D(a_4) = (x_4, x_3)$ ,  $\psi_D(a_5) = (x_4, x_2)$ ,  $\psi_D(a_6) = (x_4, x_2)$ ,  
 $\psi_D(a_7) = (x_5, x_2)$ ,  $\psi_D(a_8) = (x_2, x_5)$ ,  $\psi_D(a_9) = (x_3, x_5)$ .

$a_5, a_6$  are parallel edges,

$a_7, a_8$  are not.

$a_3$  is a loop.





# 1.1 Graph and Graphical Presentation

- Def: **graphical presentation**

- ex: in Example 1.1.1:

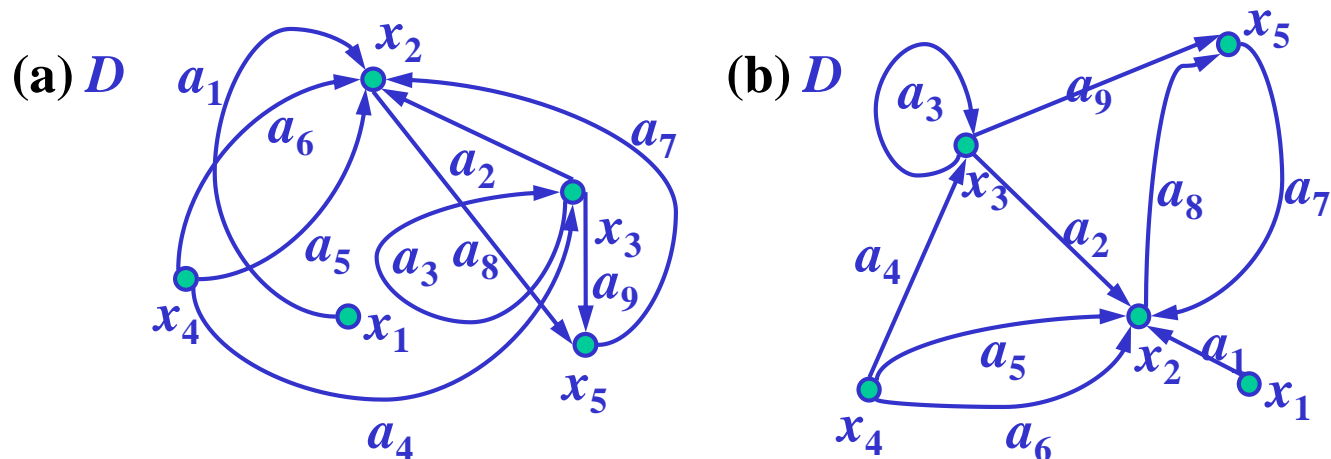
$D = (V(D), E(D), \psi_D)$  is a digraph, where

$$V(D) = \{x_1, x_2, x_3, x_4, x_5\},$$

$$E(D) = \{a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9\}$$

and  $\psi_D$  is defined by  $\psi_D(a_1) = (x_1, x_2)$ ,  $\psi_D(a_2) = (x_3, x_2)$ ,  $\psi_D(a_3) = (x_3, x_3)$ ,  
 $\psi_D(a_4) = (x_4, x_3)$ ,  $\psi_D(a_5) = (x_4, x_2)$ ,  $\psi_D(a_6) = (x_4, x_2)$ ,  
 $\psi_D(a_7) = (x_5, x_2)$ ,  $\psi_D(a_8) = (x_2, x_5)$ ,  $\psi_D(a_9) = (x_3, x_5)$ .

- **Figure 1.1,**





# 1.1 Graph and Graphical Presentation

- ex: in Example 1.1.3:

$G = (V(G), E(G), \psi_G)$  is an undirected graph, where

$$V(G) = \{z_1, z_2, z_3, z_4, z_5, z_6\},$$

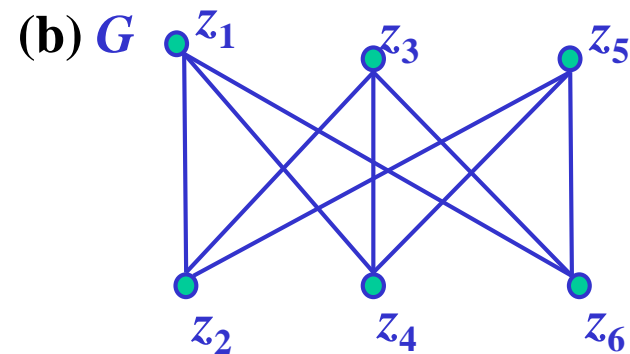
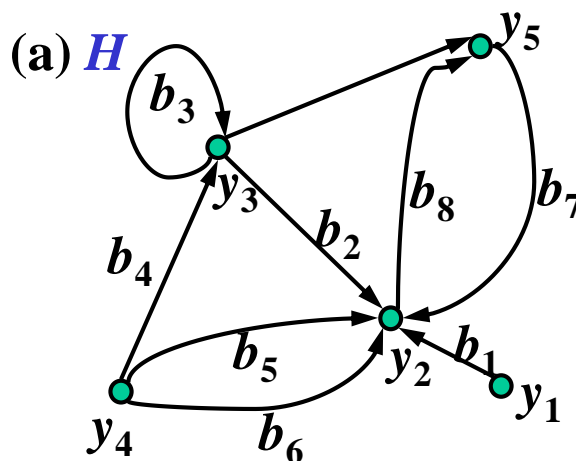
$$E(G) = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7, e_8, e_9\}$$

and  $\psi_G$  is defined by  $\psi_G(e_1) = z_1z_2$ ,  $\psi_G(e_2) = z_1z_4$ ,  $\psi_G(e_3) = z_1z_6$ ,

$$\psi_G(e_4) = z_2z_3, \psi_G(e_5) = z_3z_4, \psi_G(e_6) = z_3z_6,$$

$$\psi_G(e_7) = z_2z_5, \psi_G(e_8) = z_4z_5, \psi_G(e_9) = z_5z_6.$$

- **Figure 1.2**



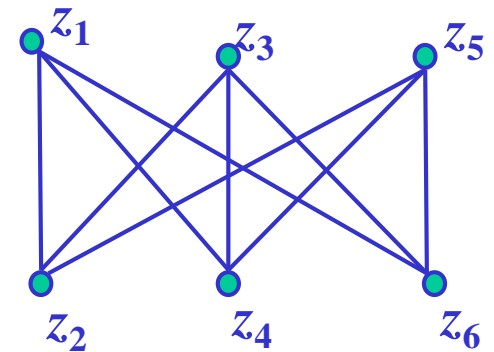


# 1.1 Graph and Graphical Presentation

- Def:
  - **incident:** vertex  $\leftrightarrow$  edge.
  - **adjacent:** vertex  $\leftrightarrow$  vertex; edge  $\leftrightarrow$  edge.
  - **loopless**
  - **simple:** a graph contains neither loops nor parallel edges.
    - $\psi$  is injective (1-1)
    - use  $V \times V$  instead of  $E$ . i.e. write  $(V, E)$  for  $(V, E, \psi)$

- **ex:** In Example 1.1.3, write as :  $G = (V(G), E(G))$ , where

$$E(G) = \{z_1z_2, z_1z_4, z_1z_6, z_2z_3, z_3z_4, z_3z_6, z_2z_5, z_4z_5, z_5z_6\}.$$





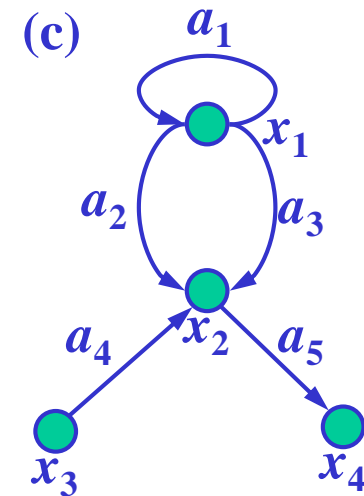
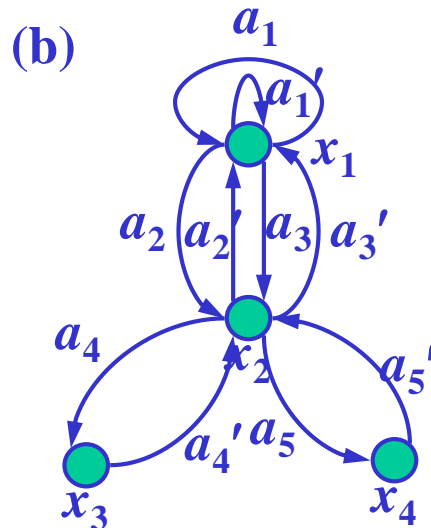
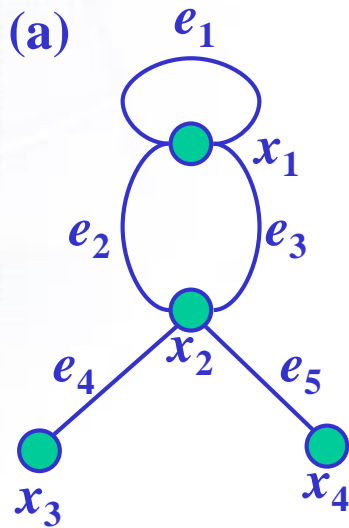
# 1.1 Graph and Graphical Presentation

- Def:
  - An undirected graph can be thought of as a **symmetric digraph**, in which there are two **symmetric edges**, one in each direction, corresponding to each undirected edge.  
→ digraph is more general !!
  - $G$  is the **underlying graph** of  $D \equiv D$  is an **oriented graph** of  $G$ .



# 1.1 Graph and Graphical Presentation

- **ex: Figure 1.3** (a) an undirected graph,  
(b) (a)'s symmetric digraph,  
(c) (a)'s an oriented graph.





# 1.1 Graph and Graphical Presentation

- Def: Let  $G = (V, E, \psi)$  be a graph,
  - $\nu = |V| (= n(G))$ , is called **order** of  $G$ . (書上印成 $\nu$ )
  - $\varepsilon = |E| (= m(G))$ , is called **size** of  $G$ .
  - A graph is **empty** if  $\varepsilon = 0$ .
  - A empty graph is **trivial** if  $\nu = 1$ .
  - **non-trivial**  $\equiv$  not trivial.
  - A graph is **finite** if both  $\nu, \varepsilon$  are finite. (本書只討論 finite graph)
  - $\lceil r \rceil, \lfloor r \rfloor, \binom{n}{k}$



# 1.1 Graph and Graphical Presentation

- **Example 1.1.4:** In any group of six people, there must be three people who get to either know each other or not.

**Proof.**

Let points  $A, B, C, D, E, F$  on the plane to denote these 6 people.

Draw **red line** joining two points if they have known each other,

**blue line** joining two points if not.

Now, consider  $F$ , there exist three lines of the same color which are incident with  $F$ .

(Without loss of generality) W.L.O.G.,

suppose they are three red lines  $FA, FB, FC$ .

Consider  $\triangle ABC$ , if it has no red line, then  $\exists$  blue  $\triangle ABC$ ,

else it has a red line, say  $AB$ , then  $\exists$  red  $\triangle ABF$ .

That means, either  $\exists$  three people ( $ABF$ ) know each other,

or  $\exists$  three people ( $ABC$ ) don't know each other.



# 1.1 Graph and Graphical Presentation

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- **Exercise: 1.1.1,** (剩下的請至少看過題目)