

Computer Science and Information Engineering
National Chi Nan University

Discrete Mathematics

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Chapter 1 Fundamental Principles of Counting

§ 1.3 Combinations: The Binomial Theorem

*Slides for a Course Based on the Text
Discrete & Combinatorial Mathematics (5th Edition)
by Ralph P. Grimaldi*

§ 1.3 Combinations: The Binomial Theorem

See textbook from p.15 to p.19 by yourself.

If there are n distinct objects, each **selection**, or **combination**, of r of these objects, with no reference to order, corresponds to $r!$ permutations of size r from the n objects. Thus the number of combinations of size r from a collection of size n , denoted $C(n, r)$, where $0 \leq r \leq n$, satisfies $(r!) \times C(n, r) = P(n, r)$ and

$$C(n, r) = \frac{P(n, r)}{r!} = \frac{n!}{r!(n-r)!}, \quad 0 \leq r \leq n.$$

Also denoted by $\binom{n}{r}$ ($= {}^n C_r$), read as “ n choose r ”.

Def : 1. For $r = 0$, $C(n, 0) = 1$, for all $n \geq 0$.

2. For all $0 \leq r \leq n$, $C(n, r) = n!/[r!(n-r)!]$.

§ 1.3 Combinations: The Binomial Theorem

EX 1.25 : Select 5 cards from 52 cards (5/52):

a) no clubs $\binom{39}{5} = 575,757$.

b) at least one club \equiv not in (a) $= \binom{52}{5} - \binom{39}{5} = 2,023,203$.

c) obtain (b) in another way ? Select one club, then 4/51 ?
 $\binom{13}{1} \cdot \binom{51}{4} = 13 \times 249,900 = 3,248,700$.

→ wrong!!

ex : 3C : 5C-KC-7H-JS

5C : 3C-KC-7H-JS

KC : 3C-5C-7H-JS

→ all the same!! But count 3 times!!

§ 1.3 Combinations: The Binomial Theorem

d) another way :

| Number of Clubs | | Number of Cards that are not Clubs | |
|-----------------|-----------------|------------------------------------|-----------------|
| 1 | $\binom{13}{1}$ | 4 | $\binom{39}{4}$ |
| 2 | $\binom{13}{2}$ | 3 | $\binom{39}{3}$ |
| 3 | $\binom{13}{3}$ | 2 | $\binom{39}{2}$ |
| 4 | $\binom{13}{4}$ | 1 | $\binom{39}{1}$ |
| 5 | $\binom{13}{5}$ | 0 | $\binom{39}{0}$ |

$$\Rightarrow \sum_{i=1}^5 \binom{13}{i} \binom{39}{5-i} = 2,023,203.$$

§ 1.3 Combinations: The Binomial Theorem

Note : $\binom{n}{r} = \binom{n}{n-r}$, $\forall n \geq r \geq 0$.

Thm 1.1 : **The Binomial Theorem**

If x and y are variables and n is a positive integer , then

$$(x + y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k} = \sum_{k=0}^n \binom{n}{n-k} x^k y^{n-k}$$

ex : $(x + y)^4 = (x + y) (x + y) (x + y) (x + y) \rightarrow x^2 y^2$

Proof.

$$(x + y)^n = (x + y) (x + y) \dots (x + y)$$

*n*th factor

the coefficient of $x^k y^{n-k}$, where $0 \leq k \leq n$, is the number of different ways in which we can select k x 's from the n available factors

§ 1.3 Combinations: The Binomial Theorem

EX 1.26 : a) $(x + y)^7$: the coefficient of $x^5 y^2 = ?$ $\binom{7}{2} = \binom{7}{5} = 21$

b) $(2a - 3b)^7$: the coefficient of $a^5 b^2 = ?$

$$\begin{aligned}(2a - 3b)^7 &= \sum_{k=0}^7 \binom{7}{k} (2a)^k (-3b)^{7-k} \\ &= \sum_{k=0}^7 \binom{7}{k} 2^k \cdot (-3)^{7-k} \cdot a^k \cdot b^{7-k}\end{aligned}$$

$$k = 5 \Rightarrow \binom{7}{5} \cdot 2^5 \cdot (-3)^{7-5} = 21 \cdot 32 \cdot 9 = 6048$$

Corollary 1.1 : $\forall n \in \mathbb{N}$, (a) $\binom{n}{0} + \binom{n}{1} + \binom{n}{2} + \dots + \binom{n}{n} = 2^n$

(b) $\binom{n}{0} - \binom{n}{1} + \binom{n}{2} - \dots + (-1)^n \binom{n}{n} = 0$

Proof.

$$(x + y)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k \cdot y^{n-k}$$

(a) Let $x = y = 1 \Rightarrow 2^n = \sum_{k=0}^n \binom{n}{k} = \binom{n}{0} + \binom{n}{1} + \binom{n}{2} + \dots + \binom{n}{n}$

(b) Let $x = -1, y = 1 \Rightarrow (-1 + 1)^n = 0 = \sum_{k=0}^n \binom{n}{k} (-1)^k (1)^{n-k}$
 $= \binom{n}{0} - \binom{n}{1} + \dots + (-1)^n \binom{n}{n}$

§ 1.3 Combinations: The Binomial Theorem

Theorem 1.2 : multinomial theorem

$$\forall n, t \in \mathbb{N}$$
$$(x_1 + x_2 + \dots + x_t)^n = \sum_{\substack{n_i \in \mathbb{Z} \ 0 \leq n_i \leq n \ \forall 1 \leq i \leq t \\ \text{and } n_1 + n_2 + \dots + n_t = n}} \frac{n!}{n_1! n_2! \dots n_t!} x_1^{n_1} x_2^{n_2} \dots x_t^{n_t}$$

Proof.

$$\begin{aligned} & \binom{n}{n_1} \binom{n-n_1}{n_2} \dots \binom{n-n_1-n_2-\dots-n_{t-1}}{n_t} \\ &= \frac{n!}{n_1!(n-n_1)!} \frac{(n-n_1)!}{n_2!(n-n_1-n_2)!} \dots \frac{(n-n_1-n_2-\dots-n_{t-1})!}{n_t!(n-n_1-n_2-\dots-n_t)!} \\ &= \frac{n!}{n_1! n_2! \dots n_t!} \end{aligned}$$

§ 1.3 Combinations: The Binomial Theorem

Def : $\frac{n!}{n_1! n_2! \cdots n_t!} = \binom{n}{n_1 n_2 \cdots n_t}$ ($= \binom{n}{n_1 n_2 \cdots n_t}$) called **multinomial coefficient**.

EX 1.27 : a) $(x + y + z)^7$:

the coefficient of $x^2 y^2 z^3 = 7! / (2! 2! 3!) = 210$

the coefficient of $x y z^5 = 7! / 5! = 42$

the coefficient of $x^3 z^4 = 7! / (3! 4!) = 35$

b) $(a + 2b - 3c + 2d + 5)^{16}$:

the coefficient of $a^2 b^3 c^2 d^5 = ?$

Sol.

General item : $\binom{16}{i j k l m} a^i (2b)^j (-3c)^k (2d)^l (5)^m$

Let $i = 2; j = 3; k = 2; l = 5$

$$\Rightarrow m = 16 - 2 - 3 - 2 - 5 = 4$$

\Rightarrow the coefficient of $a^2 b^3 c^2 d^5$

$$= \binom{16}{2 3 2 5 4} 2^3 (-3)^2 2^5 5^4 = 435,891,456,000,000$$

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Chapter 1 Fundamental Principles of Counting

§ 1.4 Combinations With Repetition

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§ 1.4 Combinations With Repetition

Note : n : distinct object, r : arrangement size $\Rightarrow n^r, \forall r \geq 0$.

EX 1.28 : 7 high school freshmen,
4 purchases : **c**heeseburger, **h**ot dog, **t**aco, **f**ish sandwich.
How many purchases are possible?

Sol.

<http://zh.wikipedia.org/wiki/%E5%A2%A8%E8%A5%BF%E5%93%A5%E5%8D%B7%E9%A5%BC>

see Table 1.6 : c, c, h, h, t, t, f \Rightarrow xx | xx | xx | x
t, t, t, t, t, f, f \Rightarrow | | xxxxx | xx
 $\Rightarrow (3 + 7)! / (7! 3!) = \binom{10}{7}$

The number of combinations of n objects taken r at a time, **with repetition**, is $C(n + r - 1, r) = \frac{(n + r - 1)!}{r!(n - 1)!} = \binom{n + r - 1}{r}$

§ 1.4 Combinations With Repetition

EX 1.29 : 20 kinds of donuts, select a dozen:

$$\binom{20+12-1}{12} = \binom{31}{12} = 141,120,525.$$

EX 1.30 : 4 vice presidents / \$1000 (each multiple of \$100)

(a) $C(4 + 10 - 1, 10) = C(13, 10) = 286$

(b) each receive at least \$100:

$$C(4 + 6 - 1, 6) = 84$$

(c) each receive at least \$100, and Mona get at least \$500:

$$\begin{aligned} & C(4 + 2 - 1, 2) \\ &= C(3 + 2 - 1, 2) + C(3 + 1 - 1, 1) + C(3 + 0 - 1, 0) \\ &= 10 \end{aligned}$$

§ 1.4 Combinations With Repetition

EX 1.31 : 7 bananas, 6 oranges / 4 children : each at least one banana

Sol.

see Table 1.7 : 1)

| | |
|------------|---------------|
| 1, 2, 3 | $b b b $ |
| 2) 1, 3, 3 | $b b b $ |

$$\Rightarrow 6! / (3! 3!) = \binom{6}{3}$$

see Table 1.8 : 1)

| | |
|---------------------|-------------------|
| 1, 2, 2, 3, 3, 4 | $o oo oo o$ |
| 3) 2, 2, 2, 3, 3, 3 | $ ooo ooo $ |

$$\Rightarrow (6 + 3)! / (6! 3!) = \binom{9}{6}$$

$$\Rightarrow C(4 + 3 - 1, 3) \cdot C(4 + 6 - 1, 6) = 1680.$$

§ 1.4 Combinations With Repetition

EX 1.32 : message : 12 different symbols + 45 spaces with at least 3 spaces between each pair of consecutive symbols

How many?

Sol.

$$45 - (3 \cdot 11) = 12$$

$$12! \cdot C(11 + 12 - 1, 12) \doteq 3.097 \cdot 10^{14}$$

EX 1.33 : All integer solutions : $x_1 + x_2 + x_3 + x_4 = 7$ where $x_i \geq 0$, $\forall 1 \leq i \leq 4$.

Sol.

$$C(4 + 7 - 1, 7) = 120.$$

§ 1.4 Combinations With Repetition

Note : The following are the equivalence:

(a) The number of integer solution of the equation

$$x_1 + x_2 + \dots + x_n = r, \quad x_i \geq 0, \quad 1 \leq i \leq n.$$

(b) The number of selections, with repetitions, of size r from a collection of size n .

(c) The number of ways r identical objects can be distributed among n distinct containers.

EX 1.34 : 10 identical white marbles (彈珠) among six distinct containers!

Sol.

$$\begin{aligned} x_1 + x_2 + \dots + x_6 &= 10 \\ C^{6+10-1}_{10} &= C^{15}_{10} \Rightarrow 3003 \end{aligned}$$

§ 1.4 Combinations With Repetition

EX 1.35 : $x_1 + x_2 + \dots + x_6 < 10, x_i \geq 0, 1 \leq i \leq 6.$

Sol.

$(\Rightarrow x_1 + x_2 + \dots + x_6 = k, k \in \mathbf{N} \text{ and } 0 \leq k \leq 9)$

$\Rightarrow x_1 + x_2 + \dots + x_6 + x_7 = 10, x_i \geq 0, 1 \leq i \leq 6, x_7 > 0$

Let $x_7 - 1 = y_7$

$\Rightarrow x_1 + x_2 + \dots + x_6 + y_7 = 9, x_i \geq 0, y_7 \geq 0$

$\Rightarrow C(7 + 9 - 1, 9) = 5005.$

EX 1.36 : $(x + y)^n$: total number of terms = ?

$\Rightarrow C(2 + n - 1, n) = n + 1.$

$(w + x + y + z)^{10}$: total number of terms = ?

$\Rightarrow C(4 + 10 - 1, 10) = 286$

§ 1.4 Combinations With Repetition

EX 1.37 : (a) compositions of 4 (partitions for the number 4) :

- | | |
|----------|------------------|
| 1) 4 | 5) 2 + 1 + 1 |
| 2) 3 + 1 | 6) 1 + 2 + 1 |
| 3) 2 + 2 | 7) 1 + 1 + 2 |
| 4) 1 + 3 | 8) 1 + 1 + 1 + 1 |

(b) the number of composition for “7”.

(i) one summand : 1

**(ii) two summands : $w_1 + w_2 = 7 \Rightarrow x_1 + x_2 = 5, x_1, x_2 \geq 0$
 $\Rightarrow \binom{2+5-1}{5} = \binom{6}{5}$**

**(iii) three summands : $y_1 + y_2 + y_3 = 7$
 $\Rightarrow z_1 + z_2 + z_3 = 4, z_1, z_2, z_3 \geq 0$
 $\Rightarrow \binom{3+4-1}{4} = \binom{6}{4}$**

⋮

(vii) seven summands $\Rightarrow \binom{7+0-1}{0} = \binom{6}{0}$

$\Rightarrow \sum_{k=0}^6 \binom{6}{k} = 2^6$ (by corollary 1.1)

§ 1.4 Combinations With Repetition

Note : for each positive integer $m \Rightarrow \sum_{k=0, m-1} \binom{m-1}{k} = 2^{m-1}$

EX 1.38 : There are $2^{12-1} = 2^{11} = 2048$ compositions of 12.

How many in those compositions where each summand is even?

Sol.

For instance, $2 + 4 + 6 = 2(1 + 2 + 3)$ $2 + 8 + 2 = 2(1 + 4 + 1)$
 $8 + 2 + 2 = 2(4 + 1 + 1)$ $6 + 6 = 2(3 + 3)$.

\Rightarrow # of compositions of 12, where each summand is even
= # of compositions of 6
= 2^{6-1}
= 2^5
= 32.

§ 1.4 Combinations With Repetition

EX 1.39 : **for $i := 1$ to 20 do**
 for $j := 1$ to i do
 for $k := 1$ to j do
 Print ($i * j + k$)
print statement executed ?

Sol.

Since $1 \leq k \leq j \leq i \leq 20$,
 $\Rightarrow ({}^{20+3-1}_3) = ({}^{22}_3) = 1540$.

Note : $r (\geq 1)$ for loops : $({}^{20+r-1}_r)$ times.

§ 1.4 Combinations With Repetition

```
EX 1.40 : counter := 0  
for  $i := 1$  to  $n$  do  
    for  $j := 1$  to  $i$  do  
        counter := counter + 1
```

Sol.

As EX 1.39, the statement executed :

$$1 \leq j \leq i \leq n, \rightarrow \binom{n+2-1}{2} = \binom{n+1}{2},$$

or

$$1 + 2 + 3 + \dots + n.$$

Therefore,

$$1 + 2 + 3 + \dots + n = \sum_{i=1}^n i = \binom{n+1}{2} = n(n+1)/2.$$

§ 1.4 Combinations With Repetition

EX 1.41 : The counter at Patti and Terri's Bar has 15 bar stools.

O O E O O O O E E E O O O E O

means 10 occupied (O) stools and 5 empty (E) stools.

Say the occupancy of the 15 stools determine 7 *runs*, as shown:

OO E OOOO EEE OOO E O
run run run run run run run

Find the total number of ways 5 E's and 10 O's can determine seven runs.

Sol. (1/2)

1. Start with E :

Let x_1 count the number of E's in the first run,
 x_2 count the number of O's in the second run,
 x_3 count the number of E's in the third run, ...,
 x_7 count the number of E's in the 7th run.

§ 1.4 Combinations With Repetition

Sol. (2/2)

1. Start with E :

$$\begin{cases} x_1 + x_3 + x_5 + x_7 = 5, & x_1, x_3, x_5, x_7 > 0; \\ x_2 + x_4 + x_6 = 10, & x_2, x_4, x_6 > 0. \end{cases}$$

$$\Rightarrow \begin{cases} y_1 + y_3 + y_5 + y_7 = 1, & y_1, y_3, y_5, y_7 \geq 0; \\ y_2 + y_4 + y_6 = 7, & y_2, y_4, y_6 \geq 0. \end{cases}$$

$$\Rightarrow \binom{4+1-1}{1} \cdot \binom{3+7-1}{7} = C(4, 1) \cdot C(9, 7) = 4 \cdot 36 = 144.$$

2. Start with O :

Let w_1 count the number of O's in the first run, ...

$$\begin{cases} w_1 + w_3 + w_5 + w_7 = 10, & w_1, w_3, w_5, w_7 > 0; \\ w_2 + w_4 + w_6 = 5, & w_2, w_4, w_6 > 0. \end{cases}$$

$$\Rightarrow \binom{4+6-1}{6} \cdot \binom{3+2-1}{2} = C(9, 6) \cdot C(4, 2) = 84 \cdot 6 = 504.$$

$$\Rightarrow 144 + 504 = 648.$$

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Chapter 2 Fundamentals of Logic

§ 2.1 Basic Connectives and Truth tables

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§ 2.1 Basic Connectives and Truth tables

Def : *statement (proposition)* : either *true* or *false*, not both.

- ex :
- ✓ p : Combinatorics is a required course for sophomores.
 - ✓ q : Margaret Mitchell wrote *Gone with the Wind*.
 - ✓ r : $2 + 3 = 5$.
 - ✗ “What a beautiful evening!”
 - ✗ “Get up and do your exercises.”
 - ✗ “The number x is an integer.”

Def : 1. *primitive statement*: No way to break them down into anything simpler.

2. 反之: *compound statement*

§ 2.1 Basic Connectives and Truth tables

Def : 1. **negation**; denoted by $\neg p$; read as “**not p**”.

ex : \perp ex \oplus p :

$\neg p$ = “Combinatorics is not a required course for sophomores.”

Def : 2. **compound statement**, using the following **logical connectives**.

a) **Conjunction**; denoted by $p \wedge q$; read as “**p and q**”.

b) **Disjunction**;

{ denoted by $p \vee q$; read as “**p (inclusive) or q**”.
{ denoted by $p \underline{\vee} q$; read as “**p exclusive or q**”.

§ 2.1 Basic Connectives and Truth tables

Def : 2. *compound statement*.

- c) *Implication*; denoted by $p \rightarrow q$; read as “*p implies q*”.
- ≡ (i) If p , then q
 - (ii) p is *sufficient* for q
 - (iii) p is a *sufficient condition* for q
 - (iv) p only if q
 - (v) q is *necessary* for p
 - (vi) q is a *necessary condition* for p
 - (vii) p is called the *hypothesis* of the implication.
 - (viii) q is called the *conclusion* of the implication.
- d) *Biconditional*; denoted by $p \leftrightarrow q$; read as “*p if and only if q*”.
- (i) “ p is *necessary and sufficient* for q .”
 - (ii) “ p *iff* q .”

§ 2.1 Basic Connectives and Truth tables

Def : *truth table*: “0” for false and “1” for true.

Table 2.1:

2.2:

| p | $\neg p$ |
|-----|----------|
| 0 | 1 |
| 1 | 0 |

| p | q | $p \wedge q$ | $p \vee q$ | $p \underline{\vee} q$ | $p \rightarrow q$ | $p \leftrightarrow q$ |
|-----|-----|--------------|------------|------------------------|-------------------|-----------------------|
| 0 | 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 | 0 | 1 | 1 |

ex : “If $2 + 3 = 6$, then $2 + 4 = 7$ ” is true.

§ 2.1 Basic Connectives and Truth tables

EX 2.1 : s : Phyllis goes out for a walk.

t : The moon is out.

u : It is snowing.

a) $(t \wedge \neg u) \rightarrow s$:

b) $t \rightarrow (\neg u \rightarrow s)$: “ $\neg u \rightarrow s$ ” means “ $(\neg u) \rightarrow s$ ”,

c) $\neg (s \leftrightarrow (u \vee t))$: not “ $\neg (u \rightarrow s)$ ”

d) “Phyllis will go out walking if and only if the moon is out” : $s \leftrightarrow t$

e) “If it is snowing and the moon is not out, then Phyllis will not go out for a walk” : $(u \wedge \neg t) \rightarrow \neg s$

f) “It is snowing but Phyllis will still go out for a walk” : $u \wedge s$
(where “but” \equiv “and”)

§ 2.1 Basic Connectives and Truth tables

EX 2.2 : “If I weigh more than 120 pounds, then I shall enroll in an exercise class”.

p: I weigh more than 120 pounds.

q: I shall enroll in an exercise class.

Penny’s statement: $p \rightarrow q$

Case 1: $p = 1$ and $q = 1$: > 120 pounds and enrolls :

Case 2: $p = 1$ and $q = 0$: > 120 pounds but not enroll :

Case 3: $p = 0$ and $q = 0$: ≤ 120 pounds and not enroll :

Case 4: $p = 0$ and $q = 1$: ≤ 120 pounds but still enroll :

§ 2.1 Basic Connectives and Truth tables

EX 2.3 : In computer science: if-then, if-then-else.

ex : $\left\{ \begin{array}{l} \text{if } x > 2 \text{ (執行時, 給定 “}x\text{” 值, 則 “}x > 2\text{” 為一 “logical statement”)} \\ \text{then } y = 2 \text{ (“executable statement”, not “logical statement”)} \\ \text{else } y = 3 \text{ (“executable statement”, not “logical statement”)} \end{array} \right.$

ex : 生活上的 “ \rightarrow ” 與 “ \leftrightarrow ”

$s \rightarrow t$: If you do your homework, then you will get to watch the baseball game.

$t \rightarrow s$: You will get to watch the baseball game only if you do your homework.

§ 2.1 Basic Connectives and Truth tables

EX 2.4 : “Margaret Mitchell wrote *Gone with the Wind*, and if $2 + 3 \neq 5$, then combinatorics is a required course for sophomores”.

$$\equiv q \wedge (\neg r \rightarrow p)$$

| p | q | r | $\neg r$ | $\neg r \rightarrow p$ | $q \wedge (\neg r \rightarrow p)$ |
|-----|-----|-----|----------|------------------------|-----------------------------------|
| 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 | 1 |

§ 2.1 Basic Connectives and Truth tables

EX 2.5 : The truth tables for ① $p \vee (q \wedge r)$; ② $(p \vee q) \wedge r$.

| p | q | r | $q \wedge r$ | $p \vee (q \wedge r)$ | $p \vee q$ | $(p \vee q) \wedge r$ |
|-----|-----|-----|--------------|-----------------------|------------|-----------------------|
| | | | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 |

∴不可以只寫 $p \vee q \wedge r$ ，需標明為 $p \vee (q \wedge r)$ 或 $(p \vee q) \wedge r$!!

§ 2.1 Basic Connectives and Truth tables

EX 2.6 : $p \rightarrow (p \vee q), p \wedge (\neg p \wedge q)$

| p | q | $p \vee q$ | $p \rightarrow (p \vee q)$ | $\neg p$ | $\neg p \wedge q$ | $p \wedge (\neg p \wedge q)$ |
|-----|-----|------------|----------------------------|----------|-------------------|------------------------------|
| 0 | 0 | 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 | 1 | 0 |
| 1 | 0 | 1 | 1 | 0 | 0 | 0 |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 |

$p \rightarrow (p \vee q)$ is true for all truth value;
 $p \wedge (\neg p \wedge q)$ is false for all truth value.

§ 2.1 Basic Connectives and Truth tables

Def 2.1 : A compound statement is called a *contradiction* (*tautology*) if it is false (true) for all truth value assignments for its component statements, denoted by F_0 (T_0).

ex : $(p_1 \wedge p_2 \wedge \dots \wedge p_n) \rightarrow q$
only need to prove:

“when $p_1 = p_2 = \dots = p_n = 1$ and q must = 1”,
then $(p_1 \wedge p_2 \wedge \dots \wedge p_n) \rightarrow q$ is a tautology
and we have a *valid argument*.

Def : Where such p_i is called *given statements (premises)*;
 q is called *conclusion*.